

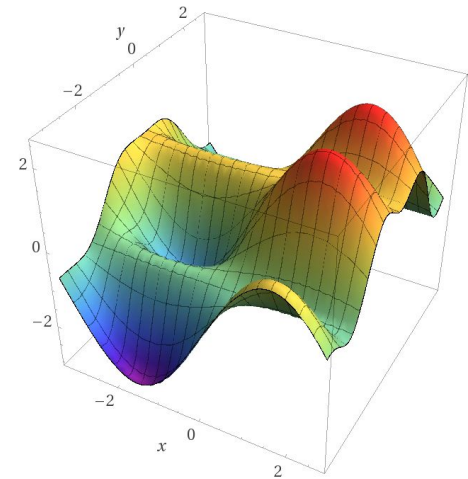
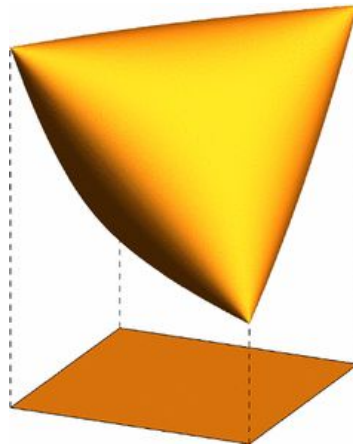
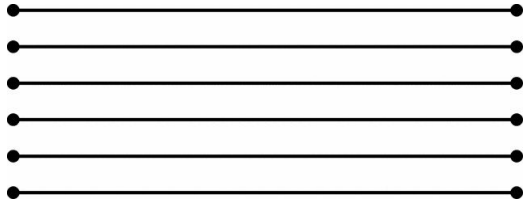
Parallel Semidefinite Optimization

An Introduction

By: Timothy Dunn



TECHNISCHE
UNIVERSITÄT
DARMSTADT



Optimization

You are a Kindergarten teacher...

... buying art supplies for students ...

... and you can't buy enough for everyone

How do you make the most kids happy?

Restrictions

- You have a budget:
 - \$24 to spend total

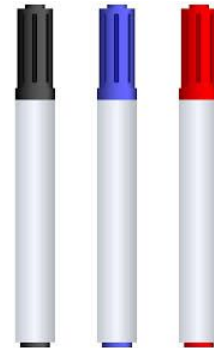
Crayons

\$1



Markers

\$2



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 - 20 students want crayons
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 - 20 students want crayons
 - 10 students want markers
- You want to be fair
 - Can't buy more than 3 more boxes of crayons than markers

How do we optimize?

How do we optimize?

let x_1 denote boxes of crayons

let x_2 denote boxes of markers

maximize $x_1 + x_2$

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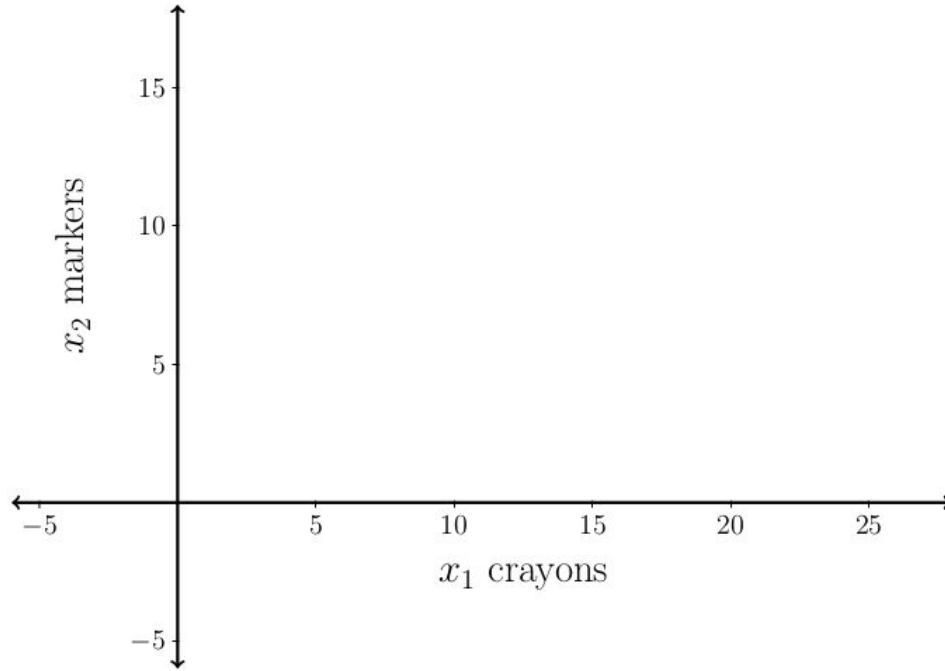
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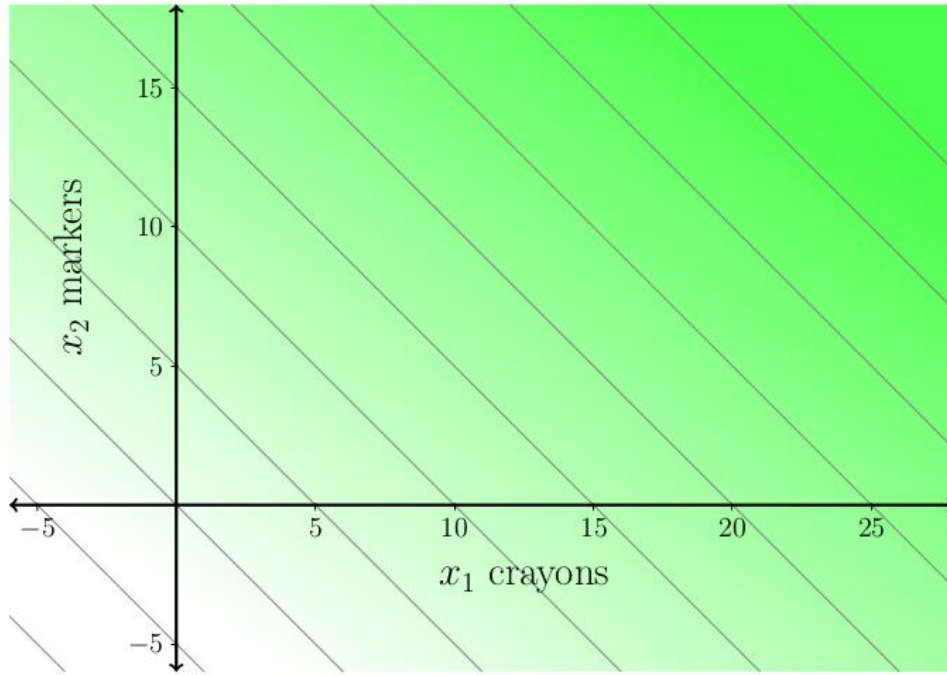
$$x_1 \leq 20 \quad x_2 \leq 10$$

$$x_1 - x_2 \leq 3$$

Optimization

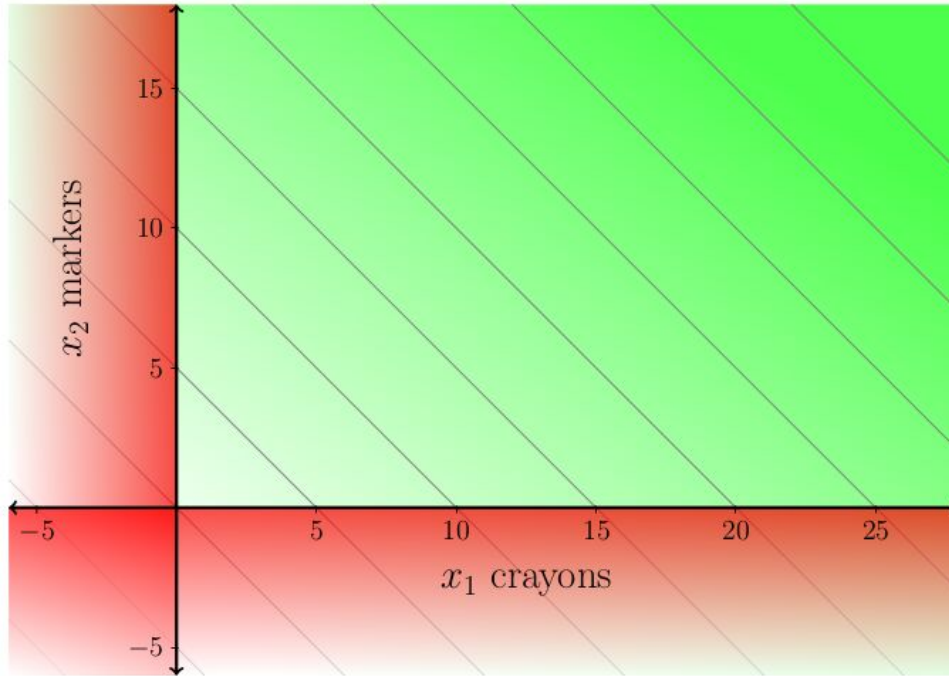


Optimization

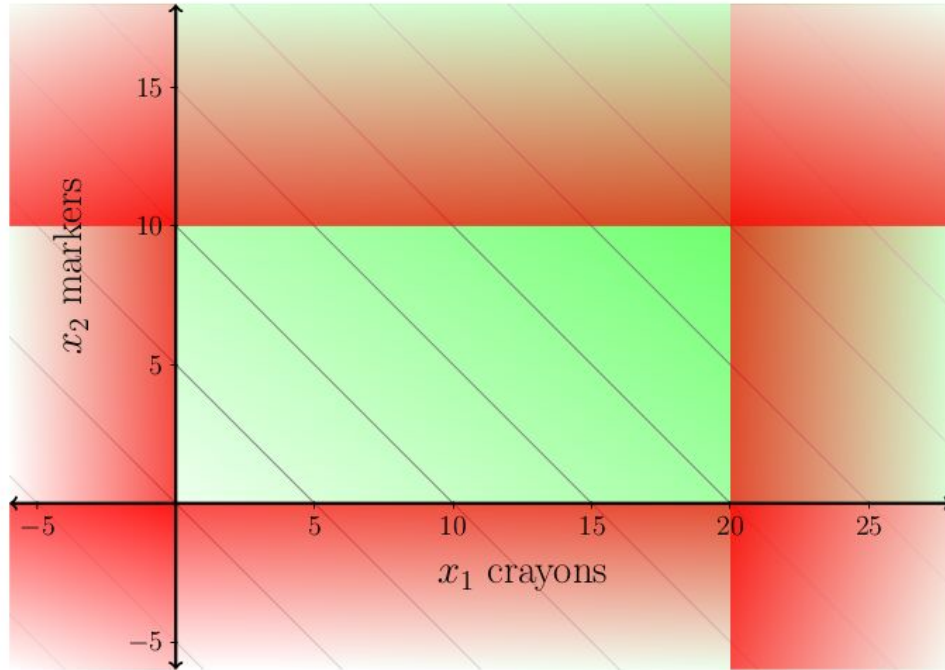


maximize
 $x_1 + x_2$

Optimization



maximize
 $x_1 + x_2$
such that
 $x_1, x_2 \geq 0$



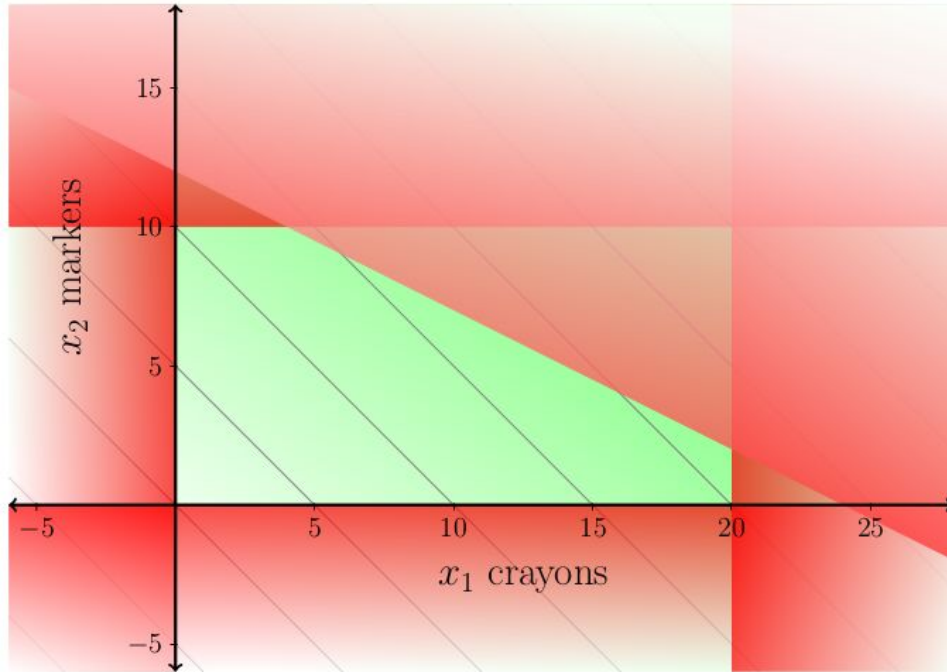
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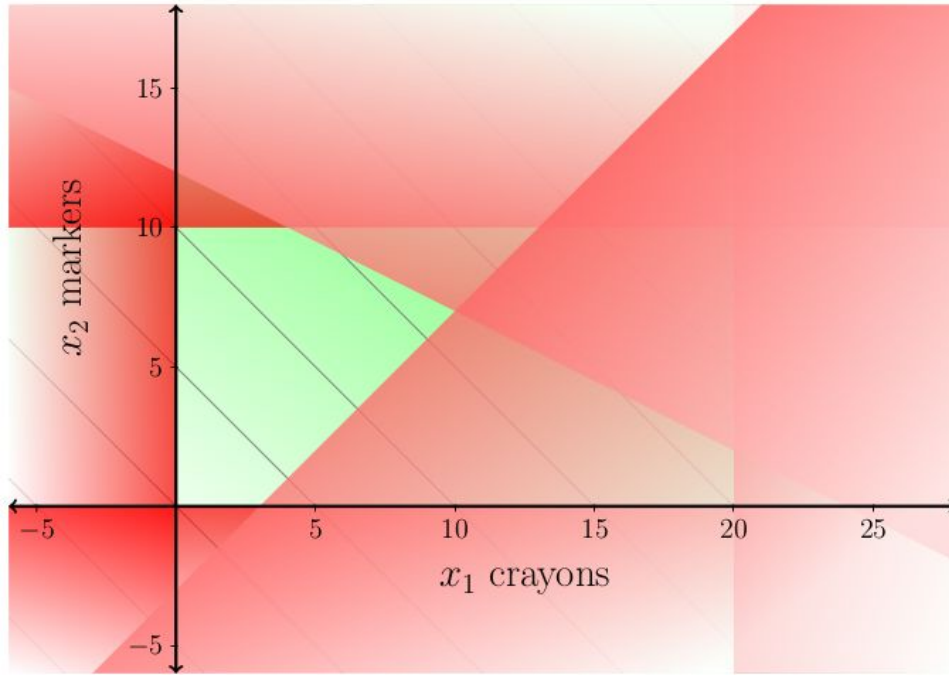
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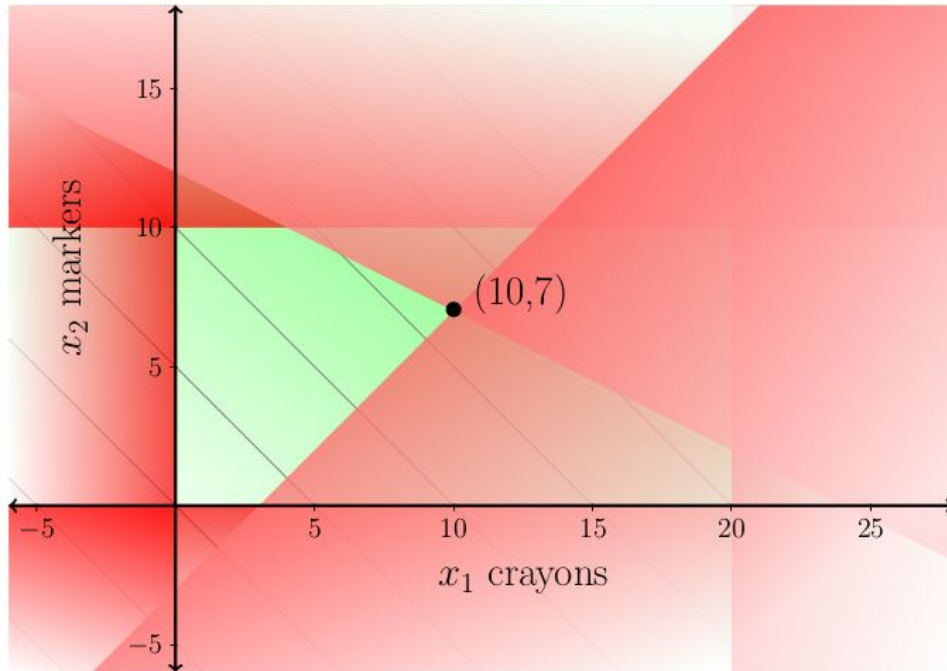
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LP Formal Definition

maximize

$$c_1x_1 + c_2x_2 + \dots + c_nx_n$$

such that

$$x_1, x_2, \dots, x_n \geq 0$$

$$a_{11}x_1 + a_{12}x_2 + \dots + a_{1n}x_n \leq b_1$$

\vdots

$$a_{m1}x_1 + a_{m2}x_2 + \dots + a_{mn}x_n \leq b_m$$

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maximize

$$\langle c^T, x \rangle$$

such that

$$x \geq 0$$

$$Ax \leq b$$

Primal

maximize

$$\langle c^T, x \rangle$$

such that

$$x \geq 0$$

$$Ax \leq b$$

Dual

minimize

$$\langle b^T, y \rangle$$

such that

$$y \geq 0$$

$$A^T y \geq c$$

LP Primal

maximize

$$\langle c^T, x \rangle$$

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SDP Primal

maximize

$$\langle C, X \rangle$$

such that

$$X \succeq 0$$

$$\langle A_k, X \rangle \leq b_k \quad \text{for } k = 1 \dots m$$

Semidefinite Programs

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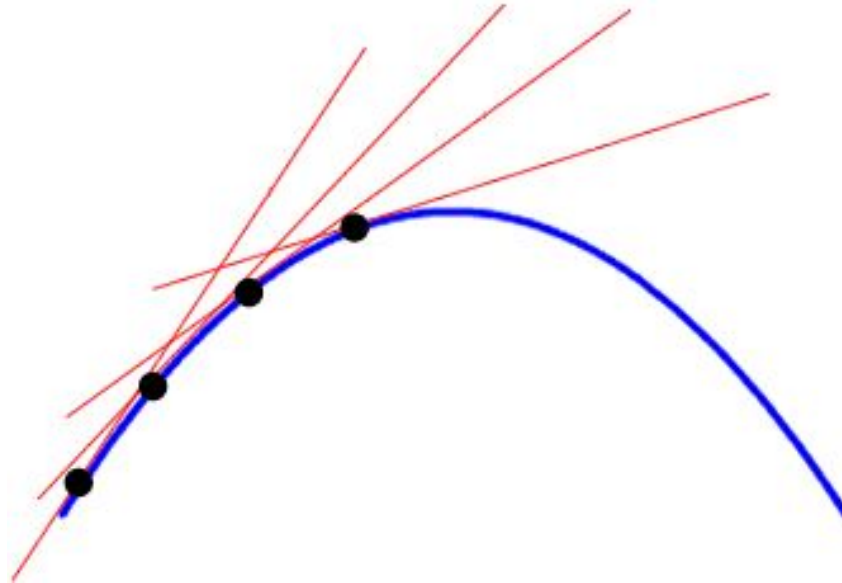
$$\forall v \in \mathbb{R}^n, \quad v^T X v \geq 0$$

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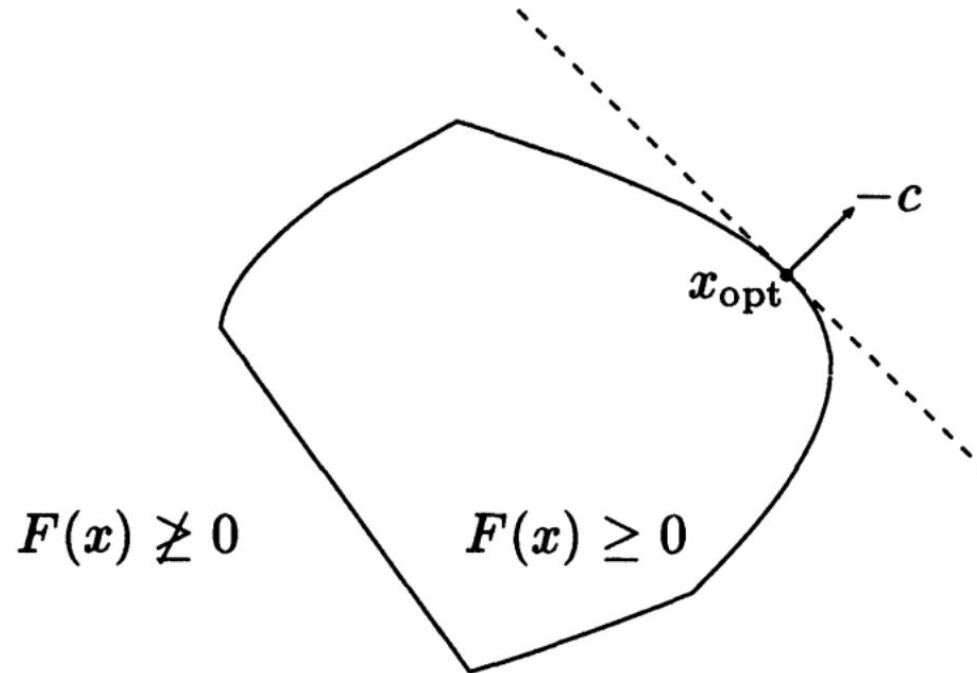
$$\forall v \in \mathbb{R}^n, \quad v^T X v \geq 0$$

This is equivalent to
an infinite set of linear constraints!

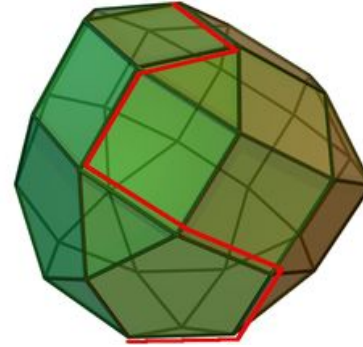
Semidefinite Programs



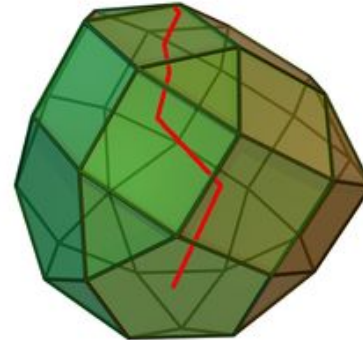
Semidefinite Programs



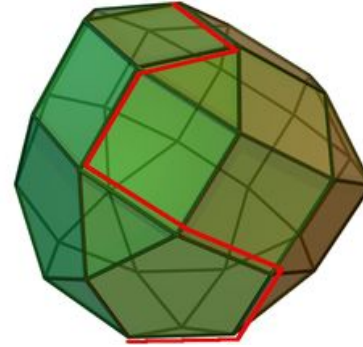
Simplex Method



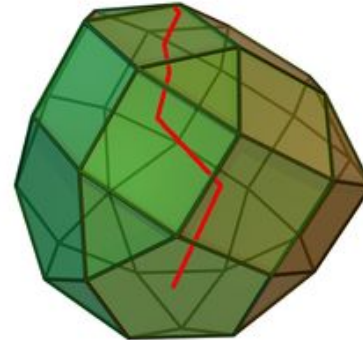
Interior Point Methods



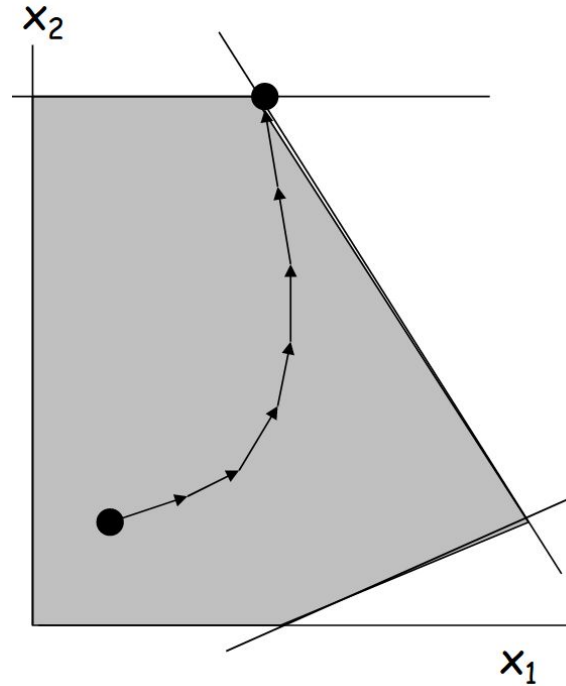
~~Simplex Method~~



Interior Point Methods



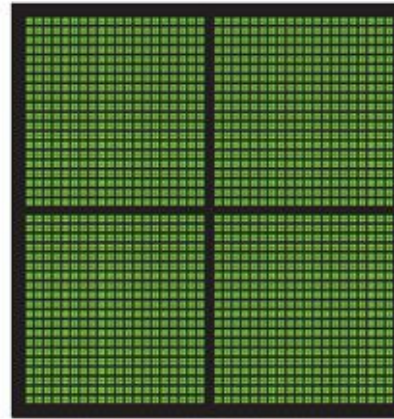
Semidefinite Programs



Parallel Programming



CPU
MULTIPLE CORES



GPU
THOUSANDS OF CORES

SDP Primal

maximize

$$\langle C, X \rangle$$

such that

$$X \succeq 0$$

$$C, X, A_k \in \mathbb{R}^{n \times n}$$

$$\langle A_k, X \rangle \leq b_k \quad \text{for } k = 1 \dots m$$

SDP Primal

maximize

$$\langle C, X \rangle$$

such that

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$$\text{let } n = 600$$

$$\text{let } m = 15$$

$$\langle A_k, X \rangle \leq b_k \quad \text{for } k = 1 \dots m$$

SDP Primal

maximize

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such that

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6 MB

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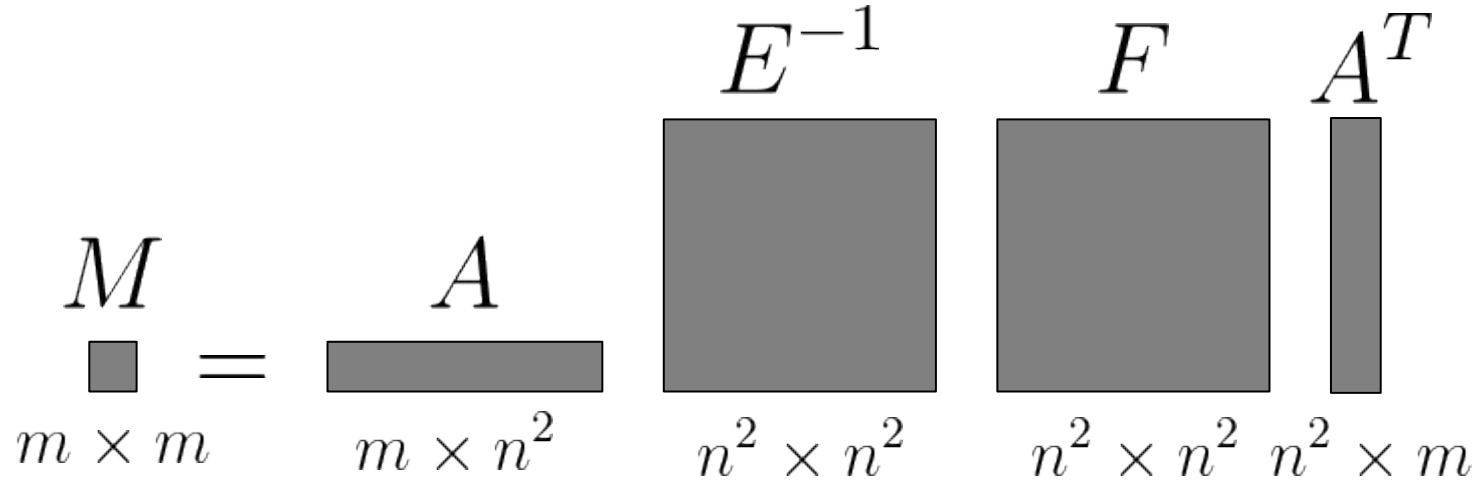
Parallel Programming: Saving Space

$$\begin{bmatrix} A & 0 & 0 \\ 0 & A^T & I \\ E & 0 & F \end{bmatrix} \begin{bmatrix} \delta x \\ \delta y \\ \delta z \end{bmatrix} = \begin{bmatrix} b - Ax \\ c - z - A^T y \\ \text{svec}(\sigma\mu I - H_P(XZ)) \end{bmatrix}$$

$$M\delta y = h$$

$$M = AE^{-1}FA^T$$

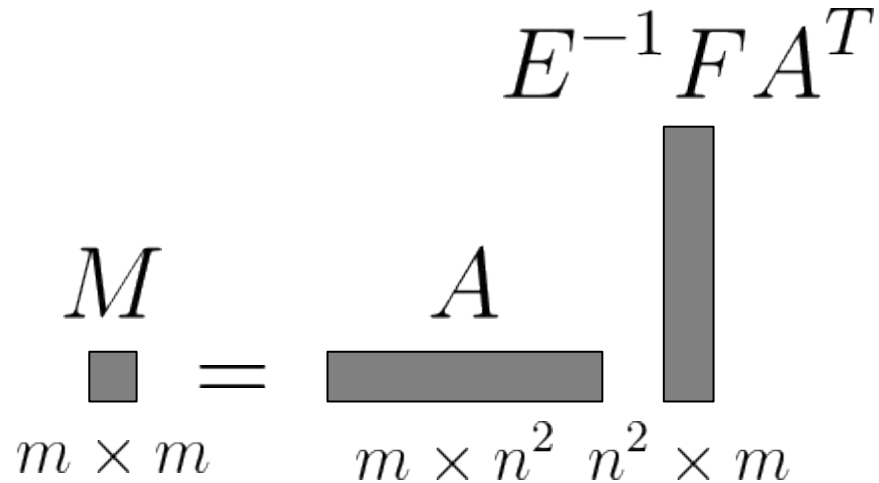
Parallel Programming: Saving Space



Parallel Programming: Saving Space

$$\begin{array}{c} M \\ \square \\ m \times m \end{array} = \begin{array}{c} A \\ \square \\ m \times n^2 \end{array} \left(\begin{array}{ccc} E^{-1} & F & A^T \\ \square & \square & \square \\ n^2 \times n^2 & n^2 \times n^2 & n^2 \times m \end{array} \right)$$

Parallel Programming: Saving Space

$$M = A E^{-1} F A^T$$


$m \times m$ $m \times n^2$ $n^2 \times n^2$ $n^2 \times m$

Parallel Programming: Saving Space

$$M = A E^{-1} F A^T$$

$m \times m$ $m \times n^2$ $n^2 \times m$

Parallel Programming: Saving Space

$$\begin{array}{c}
 M \\
 \small{m \times m}
 \end{array}
 =
 \begin{array}{c}
 A \\
 \small{m \times n^2}
 \end{array}
 \begin{array}{c}
 E^{-1} F A^T \\
 \small{n^2 \times m}
 \end{array}
 =
 \frac{1}{2}
 \begin{bmatrix}
 Z^{-1} & A_k & X \\
 X & A_k & Z^{-1}
 \end{bmatrix}$$

The diagram illustrates the decomposition of a matrix M into a product of a matrix A and a matrix $E^{-1} F A^T$. The matrix A is shown as a horizontal gray bar with a green line on top, and $E^{-1} F A^T$ is shown as a vertical gray bar with a green line on the left. Red diamonds mark the corners of these matrices. The resulting matrix is a symmetric block matrix with Z^{-1} and A_k blocks on the diagonal and X blocks off-diagonal. The A_k blocks are highlighted in green.



Questions?

Acknowledgements

Daniel Thuerck
DAAD Rise Program
Clarkson Honors Program

- [1] <https://drawinglics.com/photos/2090385/children-holding-hands-clipart-10113-kids-holding-hands-vector-art-illustration.py>
- [2] <https://clipartpng.com/?1368.markers-png-clip-art>
- [3] <https://link.springer.com/article/10.1007/s10107-015-0922-1>
- [4] <https://www.matroid.com/blog/post/the-hard-thing-about-deep-learning>
- [5] https://en.wikipedia.org/wiki/Simplex_algorithm
- [6] <http://www.stat.cmu.edu/~ryantibs/convexopt-F13/lectures/18-semidefiniteprogramming.pdf>
- [7] <http://slideplayer.com/slide/8074853/>
- [8] <https://people.duke.edu/~ccc14/sta-663/CUDAPython.html>